

Chapter 5, "Types"

Background

- Haskell's basically the *Lambda Calculus* with *types*.
- *Types* layer over **LC** to make sure evaluation doesn't *get stuck*

Summary (1/4)

- `:type 't'` - query a type in the REPL
- `(->) a b` - the *arrow* or *function* type (functions from a to b).
a and b are *type variables*.
- `(+) :: Num a => a -> a -> a` - *typeclass constrained*
type variables.
(Note: also valid are multiple constraints: `(Num a, Num b)`
`=> a -> b -> a`)

Summary (2/4)

- Arguments to functions are *curried*, or passed one-at-a-time.

$(+) :: \text{Num } a \Rightarrow a \rightarrow a \rightarrow a$

is really

$(+) :: \text{Num } a \Rightarrow a \rightarrow (a \rightarrow a)$

- Partial application

`addStuff :: Integer -> Integer -> Integer`

`addStuff a b = a + b + 5`

`let addTen = addStuff 5`

Summary (3/4)

- Explicit `curry` and `uncurry` functions

`curry` :: $((a, b) \rightarrow c) \rightarrow (a \rightarrow b \rightarrow c)$

`uncurry` :: $(a \rightarrow b \rightarrow c) \rightarrow ((a, b) \rightarrow c)$

- Operator section

`(+1)` :: $\text{Num } a \Rightarrow a \rightarrow a$

argument order matters!

`(2^)` means "2 to the power of *arg*"

`(^2)` means "*arg* squared"

Summary (4/4)

- "parametric polymorphism" - a type variable in a signature means it can be *any type*.
- $\text{id} :: a \rightarrow a$ - identity function, because of *parametricity*, we know that all this function can do is return its argument (it can't have any type-dependent behavior)
- "type inference" - Haskell determines the type of expressions automatically (when possible)

Next time

Chapter 6, "Typeclasses"

Exercise template

Make this source file compile and make the tests pass.

<https://gist.github.com/>

[twopoint718/1c46ef58ee2dbc41ea186035938e97d2](https://gist.github.com/twopoint718/1c46ef58ee2dbc41ea186035938e97d2)